



AD Master Class: Advanced Adjoint Techniques

**Checkpointing
and external functions:
Manipulating the DAG**

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AD Masterclass Schedule and Remarks

■ AD Masterclass Schedule

- ☐ 1 October 2020 | Checkpointing and external functions 1
- ☐ 15 October 2020 | Checkpointing and external functions 2
- ☐ 29 October 2020 | Guest lecture by Prof Uwe Naumann on Advanced AD topics in Machine Learning
- ☐ 12 November 2020 | Monte Carlo
- ☐ 19 November 2020 | Guest lecture by Prof Uwe Naumann on Adjoint Code Design Patterns applied to Monte Carlo
- ☐ 25 November 2020 | Computing Hessians

■ Remarks

- ☐ Please submit your questions via the questions panel at any time during this session, these will be addressed at the end.
- ☐ A recording of this session, along with the slides will be shared with you in a day or two.

We want this webinar series to be interactive (even though it's hard to do)

- We want your feedback, we want to adapt material to your feedback
- Please feel free to contact us via email to ask questions at any time
- We'd love to reach out offline, discuss what's working, what to spend more time on
- For some orgs, may make sense for us to do a few bespoke sessions

- This is an advanced course
- We assume that you are familiar with the material from the first Masterclass series
- You will get access to the materials from the first Masterclass series via email in a day or two
- Also it is not a pre-requisite we recommend to review the material from the previous series
- We will try to give references to the previous Masterclass series whenever possible

Outcomes

- How to make and fill a gap
- Difference between make/fill gap and Jacobian preaccumulation
- Use gaps to control memory usage in adjoint mode.

Algorithmic Differentiation

$$F : \mathbb{R}^n \rightarrow \mathbb{R}^m, \quad \mathbf{y} = F(\mathbf{x})$$

■ Tangent-Linear Model (TLM) \dot{F} (forward mode)

$$\dot{F}(\mathbf{x}, \dot{\mathbf{x}}) = \underset{\in \mathbb{R}^{m \times n}}{F'(\mathbf{x})} \cdot \underset{\in \mathbb{R}^n}{\dot{\mathbf{x}}} = \dot{\mathbf{y}}$$

- $F'(\mathbf{x})$ at $O(n) \cdot \text{Cost}(F)$
- exact derivatives
- $\frac{\text{Cost}(\dot{F})}{\text{Cost}(F)} \approx 2$

■ Adjoint Model (ADM) \bar{F} (reverse mode)

$$\bar{F}(\mathbf{x}, \bar{\mathbf{y}}) = \underset{\in \mathbb{R}^m}{\bar{\mathbf{y}}} \cdot \underset{\in \mathbb{R}^{m \times n}}{F'(\mathbf{x})} = F'(\mathbf{x})^T \cdot \bar{\mathbf{y}} = \bar{\mathbf{x}}$$

- $F'(\mathbf{x})$ at $O(m) \cdot \text{Cost}(F)$
- exact derivatives
- $\frac{\text{Cost}(\bar{F})}{\text{Cost}(F)} < 30$

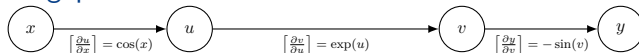
Adjoint Model

- For the adjoint model we need to reverse the control flow of the code
- AD tools record the computational graph of the program (DAG)
- Storing the DAG requires significant amount of memory.
- To run the adjoint model without running out of memory, we need to manipulate the DAG

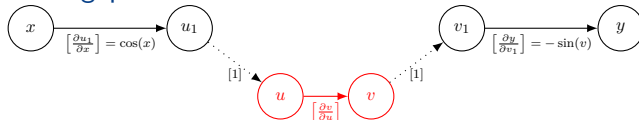
Example: Making/filling gap

$$F(x) = f \circ g \circ h(x) = \cos(\exp(\sin(x)))$$

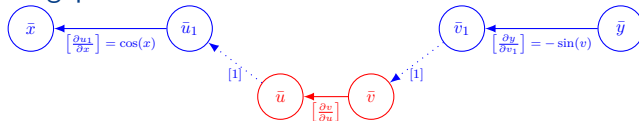
■ no gap



■ make gap



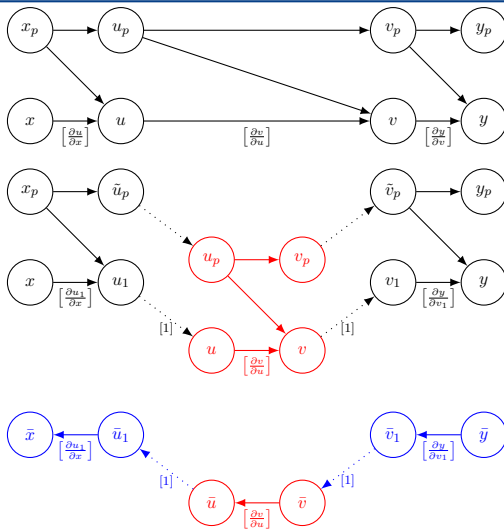
■ fill gap



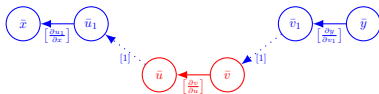
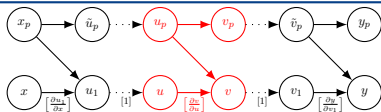
Example: Making/filling gap with AD tool

- So far we used symbolic information to fill the gap (more on this in the next masterclass)
- What if we can't differentiate the function to close the gap
- We can use the AD tool to compute the derivative and use it to fill the gap (we call this checkpointing)

Making/Filling Gap: General Case $(y, y_p) = F(x, x_p)$



Making/Filling Gap: General Case



A gap in the tape is introduced by calling a user-defined function **make_gap** to record the following **gap_data**:

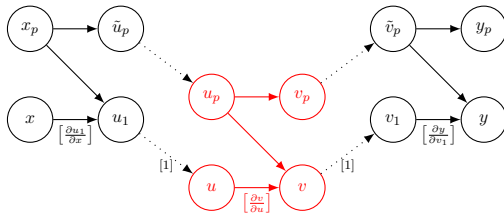
- Tape location of active gap inputs u_1 to write $\bar{u}_1 := \bar{u}$ correctly;
- adjoint gap input checkpoint $\subset (u, u_p, v, v_p)$ in order to initialize interpretation of the gap correctly;
- tape location of active gap outputs v_1 in order to initialize $\bar{v} := \bar{v}_1$ correctly; requires execution of $g(u, u_p, v, v_p)$.

This data is stored in the tape together with a reference to a user-defined function **fill_gap** to increment \bar{u}_1 with $\left(\frac{\partial v}{\partial u}\right)^T \cdot \bar{v}$.

Making Gap: Support by dco/c++

User function `make_gap`:

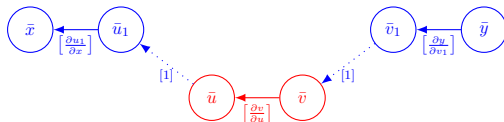
- $u := \text{dco}::\text{value}(u_1); u_p := \tilde{u}_p$
- $\text{gap_data} \rightarrow \text{write_data}(z^-)$, where $z^- \in (u_1, \tilde{u}_p, \emptyset)$
- $g(u, u_p, v, v_p)$
- $\text{dco}::\text{value}(v_1) := v$
- $\text{DCO_MODE}::\text{global_tape} \rightarrow \text{register_variable}(v_1)$
- $\text{gap_data} \rightarrow \text{write_data}(z^+)$, where $z^+ \in (v_1, v_p, \emptyset)$
- $\text{DCO_MODE}::\text{global_tape} \rightarrow \text{insert_callback}(\text{fill_gap}, \text{gap_data})$



Filling Gap (General Case): Support by dco/c++

User function `fill_gap`:

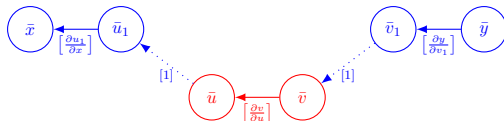
- `gap_data->read_data(z)`, where $z \equiv (z^-, z^+)$
- $\bar{v} := \text{dco}::\text{derivative}(v_1)$
- Compute the adjoint of u with $\bar{g}(z, \bar{u}, \bar{v})$
- `dco::derivative(\bar{u}_1)` + = \bar{u}



Filling Gap (Checkpointing): Support by dco/c++

User function `fill_gap`:

- `gap_data->read_data(z)`, where $z \equiv (z^-, z^+) = (u_1, \tilde{u}_p, v_1)$
- Make tape of $g(u, u_p, v, v_p)$
 - $u = u_1, u_p = \tilde{u}_p$
 - run activated version of $g(u, u_p, v, v_p)$
- set \bar{v} : `dco::derivative(v) = dco::derivative(v1)`
- Compute \bar{u} by interpreting the tape of $g(u, u_p, v, v_p)$
 - `tape->interpret_adjoint()`
- \bar{u}_1 is automatically updated with \bar{u} as $u = u_1$.



Checkpointing to reduce the required tape size

$$F(x) = f \circ g \circ h(x)$$

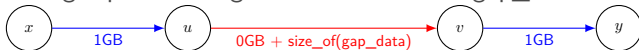
Taping F with

- no checkpoint requires 3GB for the tape

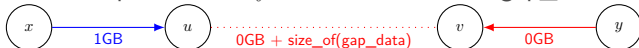


- with checkpointing of g

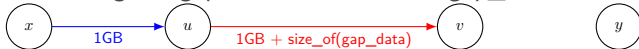
- during tape recording: 2GB + size of the gap_data



- after interpretation of f : 1GB + size of the gap_data



- after filling the gap: 2GB + size of the gap_data



Checkpointing vs. Jacobian preaccumulation

■ Jacobian preaccumulation

- ☐ performed during tape recording
- ☐ computes the Jacobian of the function, thus additional costs if the gap function has more than one output.

■ Checkpointing

- ☐ the gap is created during tape recording and filled during interpretation
- ☐ Jacobian is not computed (no additional costs if the gap function has more than one output)
- ☐ the function value of the gap function is computed twice (during make and fill gap)

Checkpointing Strategies

Developing a checkpointing strategy for code is complicated

- Checkpointing single function will typically not solve the problem of running out of memory for the adjoint mode
- A (checkpointing) strategy is necessary to avoid running out of memory for big codes
- A successful checkpointing strategy depends on many factors such as
 - ☐ size of the checkpoint (gap_data)
 - ☐ structure of the code
 - ☐ alternative ways of filling the gap

We will touch on some ideas for checkpointing strategies in the next Masterclass

Summary

- How to make/fill gap, that can be used for
 - ☐ checkpointing
 - ☐ symbolic adjoints (Masterclass 2)
 - ☐ derivatives of black box routines (Masterclass 2)
- Use make/fill gap for checkpoints (reduce memory requirements)
- Compared Jacobian preaccumulation and checkpointing

AD Master Class 2: Checkpointing and external functions 2

In the next class we will

- Learn how to use gaps to inject symbolic information into the tape for
 - ☐ linear algebra
 - ☐ root finding
 - ☐ unconstrained optimization
- Checkpointing strategies
- Look at some implementation issues in this context that need particular care e.g.
 - ☐ external function callbacks
 - ☐ memory management if code has more than one output

You will see a survey on your screen after exiting
from this session.

We would appreciate your feedback.

We are now moving on the Q&A Session